

# Combinatorial Model of Quantum Skew Symmetric Matrices

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# Quantum Points

## Definition

A point  $(x, y)$  is called Quantum if  $x$  and  $y$  satisfy the commutative relation:

$$xy = qyx$$

# Quantum Matrices

## Definition

A  $2 \times 2$  matrix is Quantum if both columns and both rows are quantum points.

$$\begin{bmatrix} a & b \\ c & d \end{bmatrix}$$

$$ab = qba$$

$$cd = qdc$$

$$ac = qca$$

$$bd = qdb$$

# Quantum Matrices

## Definition

A matrix is Quantum if all of its  $2 \times 2$  sub-matrices are quantum.

# Skew Symmetric Matrices

$$\begin{bmatrix} 0 & y_{21} & y_{31} & \dots & y_{n1} \\ -y_{21} & 0 & y_{32} & \dots & y_{n2} \\ -y_{31} & -y_{32} & 0 & \dots & y_{n3} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ -y_{n1} & -y_{n2} & -y_{n3} & \dots & 0 \end{bmatrix}$$

# The Algebra $\mathcal{O}_q(Sk_n)$

- The algebra generated by the entries in an  $n \times n$  Quantum Skew Symmetric Matrix.
- $\mathcal{O}_q(Sk_n)$  is the algebra generated by some elements

$$\{y_{ij} \mid i < j \leq n\}$$

which satisfy the relations on the next slide.

# Relations

$$y_{ij}y_{il} = qy_{il}y_{ij} \quad \text{for } i < j < l$$

$$y_{ij}y_{jl} = qy_{jl}y_{ij} \quad \text{for } i < j < l$$

$$y_{ij}y_{kj} = qy_{kj}y_{ij} \quad \text{for } i < k < j$$

$$y_{ij}y_{kl} = y_{kl}y_{ij} \quad \text{for } i < k < l < j$$

$$y_{ij}y_{kl} = y_{kl}y_{ij} + (q - q^{-1})y_{il}y_{kj} \quad \text{for } i < k < j < l$$

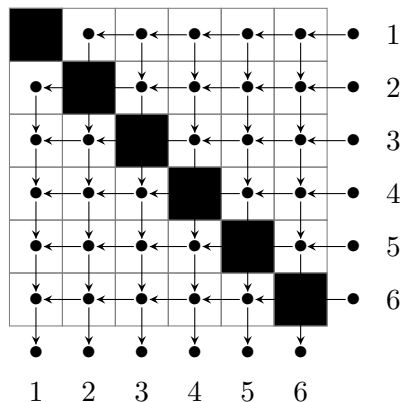
$$y_{ij}y_{kl} = y_{kl}y_{ij} + (q - q^{-1})y_{ik}y_{jl} - q(q - q^{-1})y_{il}y_{jk} \quad \text{for } i < j < k < l$$

# Goal

Find a combinatorial model for  $\mathcal{O}_q(Sk_n)$ .



## Our Model: The algebra $A_n$



- Let  $x_{ij}$  be the sum of the weights of paths from right vertex labeled  $i$  to bottom vertex labeled  $j$ .
- $A_n$  is the algebra generated by the elements of

$$\{x_{ij} \mid i < j \leq n\}$$

Figure: This graph gives  $A_6$ .

## The algebra $T_n$

Let  $T_n$  be the algebra generated by the elements of  $\{t_{ij} \mid i < j \leq n\}$  and their inverses, which satisfy the following relations:

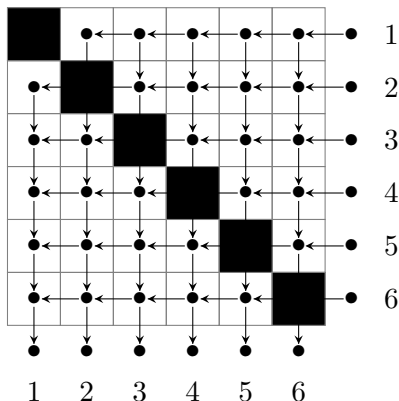
$$t_{ij}t_{il} = qt_{il}t_{ij} \quad \text{for } i < j < l$$

$$t_{ij}t_{jl} = qt_{jl}t_{ij} \quad \text{for } i < j < l$$

$$t_{ij}t_{kj} = qt_{kj}t_{ij} \quad \text{for } i < k < j$$

$$t_{ij}t_{kl} = t_{kl}t_{ij} \quad \text{for } i, j, k, l \text{ all distinct}$$

# Weighting Paths

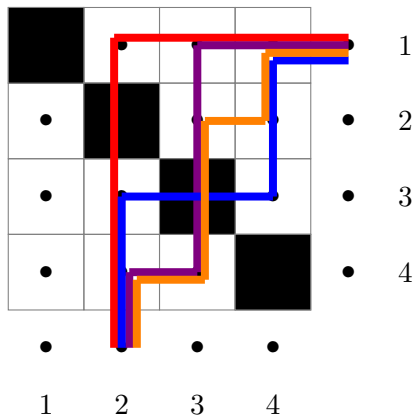


- 1** Begin with a weight of 1, and then consider the turns of the path in order.
- 2** If the path has a  $\Gamma$  turn at the  $(i, j)$ th vertex multiply by  $t_{ij}$ .
- 3** For a  $\lrcorner$  turn at the  $(i, j)$ th vertex multiply by  $t_{ij}^{-1}$ .
- 4** If the path crosses the diagonal horizontally multiply by  $-q$ .

## Example

In  $A_4$ , here is the sum of weights of paths from 1 to 2:

$$x_{12} = t_{12} + t_{13}t_{34}^{-1}t_{24} + t_{14}t_{24}^{-1}t_{23}t_{34}^{-1}t_{24} - qt_{14}t_{34}^{-1}t_{23}$$



## Theorem

$$\mathcal{O}_q(Sk_n) \cong A_n$$

- We define the map  $\phi : \mathcal{O}_q(Sk_n) \rightarrow A_n$  to be the homomorphism which sends each  $y_{ij}$  to the corresponding  $x_{ij}$ .
- Surjectivity is given by the fact that each generator is sent to.
- To show that the map is well defined, we need to show that the  $x_{ij}$ 's satisfy the commutativity relations.
- Given well defined we can use techniques from GK-dimension theory to show injectivity.

# Method

Induction on  $n$ .

Find a relation between  $x_{ij}$ 's in  $A_n$  and those in  $A_{n-1}$ .

# Parent Paths

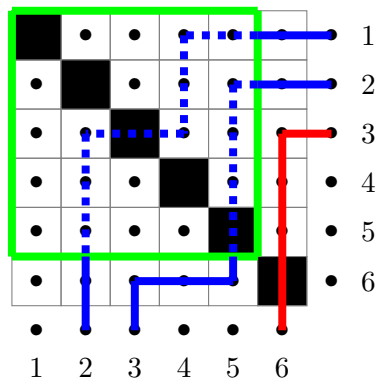
## Definition

A path  $P_{ij}$  in an  $n \times n$  graph has the *parent path*  $P_{kl}$  in an  $(n-1) \times (n-1)$  graph, if  $P_{ij}$  has the form:

$$P_{ij} = t_{in} t_{kn}^{-1} P_{kl} t_{ln}^{-1} t_{jn}$$

$P_{ij}$  is then considered a *child path* of  $P_{kl}$ .

# Parent Paths Graphically



■  $-qt_{16}t_{16}^{-1}t_{14}t_{34}^{-1}t_{23}t_{26}^{-1}t_{26}$

■  $t_{26}t_{26}^{-1}t_{25}t_{56}^{-1}t_{36}$

■  $t_{36}$



## Parental $x$ 's

$$\begin{aligned}x_{ij} &= \sum_{P_{i \rightarrow j}} w(P_{ij}) \\&= \sum_{k,l} \sum_{P_{k \rightarrow l}} t_{in} t_{kn}^{-1} w(P_{kl}) t_{ln}^{-1} t_{jn} \\&= \sum_{k,l} t_{in} t_{kn}^{-1} \left( \sum_{P_{k \rightarrow l}} w(P_{kl}) \right) t_{ln}^{-1} t_{jn} \\&= \sum_{k,l} t_{in} t_{kn}^{-1} x_{kl} t_{ln}^{-1} t_{jn}\end{aligned}$$

$$i \leq k < j \leq l < n$$

## Standard Form

$$t_{an} t_{a'n}^{-1} x_{a'b'} t_{b'n}^{-1} t_{bn} t_{cn} t_{c'n}^{-1} x_{c'd'} t_{d'n}^{-1} t_{dn}$$

## Standard Form

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## Standard Form

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- $x$ 's next to each other allows use of inductive hypothesis.
- Simplest rules for tracking  $q$  creation, based on the relations between indices.
- Terms from  $x_{cd}x_{ab}$  can be placed in standard form without creating any  $q$ 's.

# Term Cancellation

$$a \leq a' < b = c \leq c' < d < b' = d'$$

- $t_{an} t_{a'n}^{-1} x_{a'b'} t_{b'n}^{-1} t_{bn} t_{cn} t_{c'n}^{-1} x_{c'd'} t_{d'n}^{-1} t_{dn}$

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- $qt_{cn}t_{c'n}^{-1}t_{dn}t_{d'n}^{-1}x_{c'd'}x_{a'b'}t_{a'n}^{-1}t_{an}t_{b'n}^{-1}t_{bn}$

# Pair Term

$$a < c \leq a' = c' < b = d \leq b'_1 < d'_1$$



## Pair Term

$$a < c \leq a' = c' < b = d \leq b'_1 < d'_1$$

$$a < c \leq a' = c' < b = d \leq d'_2 < b'_2$$

$$b'_1 = d'_2 < b'_2 = d'_1$$

## Pair Term Cancellation

$$a < c \leq a' = c' < b = d \leq b'_1 < d'_1$$

$$(q^3 - q)t_{cn}t_{c'n}^{-1}t_{dn}t_{d'_1n}^{-1}x_{c'd'_1}x_{a'b'_1}t_{a'n}^{-1}t_{an}t_{b'_1n}^{-1}t_{bn}$$

$$a < c \leq a' = c' < b = d \leq d'_2 < b'_2$$

$$(q - q^3)t_{cn}t_{c'n}^{-1}t_{dn}t_{d'_1n}^{-1}x_{c'd'_1}x_{a'b'_1}t_{a'n}^{-1}t_{an}t_{b'_1n}^{-1}t_{bn}$$

# Cancellation Recap

Every possible ordering guarantees the existence of one of the following:

- An equal term on the opposite side of the equality.
- A pair term which, together with the original, becomes equal to two guaranteed terms on the opposite side of the equality.
- Extended versions of the second.
- Worst Case: Four terms on the left combine to equal eight on the right.

# Relations

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$$\mathcal{O}_q(Sk_n) \cong A_n$$

# Acknowledgments

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